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2's Complement Circuit

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Summary. This article introduces various Boolean operators which are used in discussing the properties and stability of a 2's complement circuit. We present the definitions and related theorems for the following logical operators which include negative input/output: 'and2a', 'or2a', 'xor2a' and 'nand2a', 'nor2a', etc. We formalize the concept of a 2's complement circuit, define the structures of complementors/incrementors for binary operations, and prove the stability of the circuit.

MML Identifier: TWOSCOMP.

WWW: http://mizar.org/JFM/Vol8/twoscomp.html

The articles [9], [11], [12], [2], [3], [13], [4], [8], [10], [6], [7], [5], and [1] provide the notation and terminology for this paper.

1. Boolean Operators

Let S be an unsplit non void non empty many sorted signature, let A be a Boolean circuit of S, let s be a state of A, and let v be a vertex of S. Then $s(v)$ is an element of *Boolean*.

The function and_2 from $Boolean^2$ into *Boolean* is defined by:

(Def. 1) For all elements x, y of *Boolean* holds $\mathrm{and}_2(\langle x,y\rangle) = x \wedge y$.

The function and_{2a} from $Boolean^2$ into *Boolean* is defined as follows:

(Def. 2) For all elements x, y of *Boolean* holds $(\mathrm{and}_{2a})(\langle x,y\rangle) = \neg x \wedge y$.

The function and_{2b} from $Boolean^2$ into *Boolean* is defined as follows:

(Def. 3) For all elements x, y of *Boolean* holds $(\mathrm{and}_{2b})(\langle x,y\rangle) = \neg x \wedge \neg y$.

The function nand_2 from $Boolean^2$ into *Boolean* is defined by:

(Def. 4) For all elements x, y of *Boolean* holds $\mathrm{nand}_2(\langle x,y\rangle) = \neg(x \wedge y)$.

The function nand_{2a} from $Boolean^2$ into *Boolean* is defined by:

(Def. 5) For all elements x, y of *Boolean* holds $(\mathrm{nand}_{2a})(\langle x,y\rangle) = \neg(\neg x \wedge y)$.

The function nand_{2b} from $Boolean^2$ into *Boolean* is defined by:

(Def. 6) For all elements x, y of *Boolean* holds $(\mathrm{nand}_{2b})(\langle x,y\rangle) = \neg(\neg x \wedge \neg y)$.

The function or_2 from $Boolean^2$ into *Boolean* is defined as follows:

(Def. 7) For all elements x, y of *Boolean* holds $\mathrm{or}_2(\langle x,y\rangle) = x \vee y$.

The function or_{2a} from $Boolean^2$ into $Boolean$ is defined by:

(Def. 8) For all elements x, y of $Boolean$ holds $(\mathrm{or}_{2a})(\langle x,y\rangle) = \neg x \vee y$.

The function or_{2b} from $Boolean^2$ into $Boolean$ is defined by:

(Def. 9) For all elements x, y of $Boolean$ holds $(\mathrm{or}_{2b})(\langle x,y\rangle) = \neg x \vee \neg y$.

The function nor_2 from $Boolean^2$ into $Boolean$ is defined by:

(Def. 10) For all elements x, y of $Boolean$ holds $\mathrm{nor}_2(\langle x,y\rangle) = \neg(x \vee y)$.

The function nor_{2a} from $Boolean^2$ into $Boolean$ is defined by:

(Def. 11) For all elements x, y of $Boolean$ holds $(\mathrm{nor}_{2a})(\langle x,y\rangle) = \neg(\neg x \vee y)$.

The function nor_{2b} from $Boolean^2$ into $Boolean$ is defined as follows:

(Def. 12) For all elements x, y of $Boolean$ holds $(\mathrm{nor}_{2b})(\langle x,y\rangle) = \neg(\neg x \vee \neg y)$.

The function xor_2 from $Boolean^2$ into $Boolean$ is defined as follows:

(Def. 13) For all elements x, y of $Boolean$ holds $\mathrm{xor}_2(\langle x,y\rangle) = x \oplus y$.

The function xor_{2a} from $Boolean^2$ into $Boolean$ is defined as follows:

(Def. 14) For all elements x, y of $Boolean$ holds $(\mathrm{xor}_{2a})(\langle x,y\rangle) = \neg x \oplus y$.

The function xor_{2b} from $Boolean^2$ into $Boolean$ is defined by:

(Def. 15) For all elements x, y of $Boolean$ holds $(\mathrm{xor}_{2b})(\langle x,y\rangle) = \neg x \oplus \neg y$.

The following propositions are true:

(3)[1] For all elements x, y of $Boolean$ holds $\mathrm{and}_2(\langle x,y\rangle) = x \wedge y$ and $(\mathrm{and}_{2a})(\langle x,y\rangle) = \neg x \wedge y$ and $(\mathrm{and}_{2b})(\langle x,y\rangle) = \neg x \wedge \neg y$.

(4) For all elements x, y of $Boolean$ holds $\mathrm{nand}_2(\langle x,y\rangle) = \neg(x \wedge y)$ and $(\mathrm{nand}_{2a})(\langle x,y\rangle) = \neg(\neg x \wedge y)$ and $(\mathrm{nand}_{2b})(\langle x,y\rangle) = \neg(\neg x \wedge \neg y)$.

(5) For all elements x, y of $Boolean$ holds $\mathrm{or}_2(\langle x,y\rangle) = x \vee y$ and $(\mathrm{or}_{2a})(\langle x,y\rangle) = \neg x \vee y$ and $(\mathrm{or}_{2b})(\langle x,y\rangle) = \neg x \vee \neg y$.

(6) For all elements x, y of $Boolean$ holds $\mathrm{nor}_2(\langle x,y\rangle) = \neg(x \vee y)$ and $(\mathrm{nor}_{2a})(\langle x,y\rangle) = \neg(\neg x \vee y)$ and $(\mathrm{nor}_{2b})(\langle x,y\rangle) = \neg(\neg x \vee \neg y)$.

(7) For all elements x, y of $Boolean$ holds $\mathrm{xor}_2(\langle x,y\rangle) = x \oplus y$ and $(\mathrm{xor}_{2a})(\langle x,y\rangle) = \neg x \oplus y$ and $(\mathrm{xor}_{2b})(\langle x,y\rangle) = \neg x \oplus \neg y$.

(8) For all elements x, y of $Boolean$ holds $\mathrm{and}_2(\langle x,y\rangle) = (\mathrm{nor}_{2b})(\langle x,y\rangle)$ and $(\mathrm{and}_{2a})(\langle x,y\rangle) = (\mathrm{nor}_{2a})(\langle y,x\rangle)$ and $(\mathrm{and}_{2b})(\langle x,y\rangle) = \mathrm{nor}_2(\langle x,y\rangle)$.

(9) For all elements x, y of $Boolean$ holds $\mathrm{or}_2(\langle x,y\rangle) = (\mathrm{nand}_{2b})(\langle x,y\rangle)$ and $(\mathrm{or}_{2a})(\langle x,y\rangle) = (\mathrm{nand}_{2a})(\langle y,x\rangle)$ and $(\mathrm{or}_{2b})(\langle x,y\rangle) = \mathrm{nand}_2(\langle x,y\rangle)$.

(10) For all elements x, y of $Boolean$ holds $(\mathrm{xor}_{2b})(\langle x,y\rangle) = \mathrm{xor}_2(\langle x,y\rangle)$.

(11) $\mathrm{and}_2(\langle 0,0\rangle) = 0$ and $\mathrm{and}_2(\langle 0,1\rangle) = 0$ and $\mathrm{and}_2(\langle 1,0\rangle) = 0$ and $\mathrm{and}_2(\langle 1,1\rangle) = 1$ and $(\mathrm{and}_{2a})(\langle 0,0\rangle) = 0$ and $(\mathrm{and}_{2a})(\langle 0,1\rangle) = 1$ and $(\mathrm{and}_{2a})(\langle 1,0\rangle) = 0$ and $(\mathrm{and}_{2a})(\langle 1,1\rangle) = 0$ and $(\mathrm{and}_{2b})(\langle 0,0\rangle) = 1$ and $(\mathrm{and}_{2b})(\langle 0,1\rangle) = 0$ and $(\mathrm{and}_{2b})(\langle 1,0\rangle) = 0$ and $(\mathrm{and}_{2b})(\langle 1,1\rangle) = 0$.

[1] The propositions (1) and (2) have been removed.

(12) $\mathrm{or}_2(\langle 0,0\rangle) = 0$ and $\mathrm{or}_2(\langle 0,1\rangle) = 1$ and $\mathrm{or}_2(\langle 1,0\rangle) = 1$ and $\mathrm{or}_2(\langle 1,1\rangle) = 1$ and $(\mathrm{or}_{2a})(\langle 0,0\rangle) = 1$ and $(\mathrm{or}_{2a})(\langle 0,1\rangle) = 1$ and $(\mathrm{or}_{2a})(\langle 1,0\rangle) = 0$ and $(\mathrm{or}_{2a})(\langle 1,1\rangle) = 1$ and $(\mathrm{or}_{2b})(\langle 0,0\rangle) = 1$ and $(\mathrm{or}_{2b})(\langle 0,1\rangle) = 1$ and $(\mathrm{or}_{2b})(\langle 1,0\rangle) = 1$ and $(\mathrm{or}_{2b})(\langle 1,1\rangle) = 0$.

(13) $\mathrm{xor}_2(\langle 0,0\rangle) = 0$ and $\mathrm{xor}_2(\langle 0,1\rangle) = 1$ and $\mathrm{xor}_2(\langle 1,0\rangle) = 1$ and $\mathrm{xor}_2(\langle 1,1\rangle) = 0$ and $(\mathrm{xor}_{2a})(\langle 0,0\rangle) = 1$ and $(\mathrm{xor}_{2a})(\langle 0,1\rangle) = 0$ and $(\mathrm{xor}_{2a})(\langle 1,0\rangle) = 0$ and $(\mathrm{xor}_{2a})(\langle 1,1\rangle) = 1$.

The function and_3 from *Boolean*^{3} into *Boolean* is defined as follows:

(Def. 16) For all elements x, y, z of *Boolean* holds $\mathrm{and}_3(\langle x,y,z\rangle) = x \wedge y \wedge z$.

The function and_{3a} from *Boolean*^{3} into *Boolean* is defined by:

(Def. 17) For all elements x, y, z of *Boolean* holds $(\mathrm{and}_{3a})(\langle x,y,z\rangle) = \neg x \wedge y \wedge z$.

The function and_{3b} from *Boolean*^{3} into *Boolean* is defined by:

(Def. 18) For all elements x, y, z of *Boolean* holds $(\mathrm{and}_{3b})(\langle x,y,z\rangle) = \neg x \wedge \neg y \wedge z$.

The function and_{3c} from *Boolean*^{3} into *Boolean* is defined as follows:

(Def. 19) For all elements x, y, z of *Boolean* holds $(\mathrm{and}_{3c})(\langle x,y,z\rangle) = \neg x \wedge \neg y \wedge \neg z$.

The function nand_3 from *Boolean*^{3} into *Boolean* is defined as follows:

(Def. 20) For all elements x, y, z of *Boolean* holds $\mathrm{nand}_3(\langle x,y,z\rangle) = \neg(x \wedge y \wedge z)$.

The function nand_{3a} from *Boolean*^{3} into *Boolean* is defined as follows:

(Def. 21) For all elements x, y, z of *Boolean* holds $(\mathrm{nand}_{3a})(\langle x,y,z\rangle) = \neg(\neg x \wedge y \wedge z)$.

The function nand_{3b} from *Boolean*^{3} into *Boolean* is defined by:

(Def. 22) For all elements x, y, z of *Boolean* holds $(\mathrm{nand}_{3b})(\langle x,y,z\rangle) = \neg(\neg x \wedge \neg y \wedge z)$.

The function nand_{3c} from *Boolean*^{3} into *Boolean* is defined by:

(Def. 23) For all elements x, y, z of *Boolean* holds $(\mathrm{nand}_{3c})(\langle x,y,z\rangle) = \neg(\neg x \wedge \neg y \wedge \neg z)$.

The function or_3 from *Boolean*^{3} into *Boolean* is defined by:

(Def. 24) For all elements x, y, z of *Boolean* holds $\mathrm{or}_3(\langle x,y,z\rangle) = x \vee y \vee z$.

The function or_{3a} from *Boolean*^{3} into *Boolean* is defined as follows:

(Def. 25) For all elements x, y, z of *Boolean* holds $(\mathrm{or}_{3a})(\langle x,y,z\rangle) = \neg x \vee y \vee z$.

The function or_{3b} from *Boolean*^{3} into *Boolean* is defined by:

(Def. 26) For all elements x, y, z of *Boolean* holds $(\mathrm{or}_{3b})(\langle x,y,z\rangle) = \neg x \vee \neg y \vee z$.

The function or_{3c} from *Boolean*^{3} into *Boolean* is defined by:

(Def. 27) For all elements x, y, z of *Boolean* holds $(\mathrm{or}_{3c})(\langle x,y,z\rangle) = \neg x \vee \neg y \vee \neg z$.

The function nor_3 from *Boolean*^{3} into *Boolean* is defined as follows:

(Def. 28) For all elements x, y, z of *Boolean* holds $\mathrm{nor}_3(\langle x,y,z\rangle) = \neg(x \vee y \vee z)$.

The function nor_{3a} from *Boolean*^{3} into *Boolean* is defined by:

(Def. 29) For all elements x, y, z of *Boolean* holds $(\mathrm{nor}_{3a})(\langle x,y,z\rangle) = \neg(\neg x \vee y \vee z)$.

The function nor_{3b} from *Boolean*^{3} into *Boolean* is defined as follows:

(Def. 30) For all elements x, y, z of *Boolean* holds $(\mathrm{nor}_{3b})(\langle x,y,z\rangle) = \neg(\neg x \vee \neg y \vee z)$.

The function nor_{3c} from *Boolean*^{3} into *Boolean* is defined as follows:

(Def. 31) For all elements x, y, z of *Boolean* holds $(\mathrm{nor}_{3c})(\langle x,y,z\rangle) = \neg(\neg x \vee \neg y \vee \neg z)$.

The function xor_3 from *Boolean*^{3} into *Boolean* is defined by:

(Def. 32) For all elements x, y, z of *Boolean* holds $\mathrm{xor}_3(\langle x,y,z\rangle) = x \oplus y \oplus z$.

The following propositions are true:

(14) For all elements x, y, z of *Boolean* holds $\mathrm{and}_3(\langle x,y,z\rangle) = x \wedge y \wedge z$ and $(\mathrm{and}_{3a})(\langle x,y,z\rangle) = \neg x \wedge y \wedge z$ and $(\mathrm{and}_{3b})(\langle x,y,z\rangle) = \neg x \wedge \neg y \wedge z$ and $(\mathrm{and}_{3c})(\langle x,y,z\rangle) = \neg x \wedge \neg y \wedge \neg z$.

(15) Let x, y, z be elements of *Boolean*. Then $\mathrm{nand}_3(\langle x,y,z\rangle) = \neg(x \wedge y \wedge z)$ and $(\mathrm{nand}_{3a})(\langle x,y,z\rangle) = \neg(\neg x \wedge y \wedge z)$ and $(\mathrm{nand}_{3b})(\langle x,y,z\rangle) = \neg(\neg x \wedge \neg y \wedge z)$ and $(\mathrm{nand}_{3c})(\langle x,y,z\rangle) = \neg(\neg x \wedge \neg y \wedge \neg z)$.

(16) For all elements x, y, z of *Boolean* holds $\mathrm{or}_3(\langle x,y,z\rangle) = x \vee y \vee z$ and $(\mathrm{or}_{3a})(\langle x,y,z\rangle) = \neg x \vee y \vee z$ and $(\mathrm{or}_{3b})(\langle x,y,z\rangle) = \neg x \vee \neg y \vee z$ and $(\mathrm{or}_{3c})(\langle x,y,z\rangle) = \neg x \vee \neg y \vee \neg z$.

(17) Let x, y, z be elements of *Boolean*. Then $\mathrm{nor}_3(\langle x,y,z\rangle) = \neg(x \vee y \vee z)$ and $(\mathrm{nor}_{3a})(\langle x,y,z\rangle) = \neg(\neg x \vee y \vee z)$ and $(\mathrm{nor}_{3b})(\langle x,y,z\rangle) = \neg(\neg x \vee \neg y \vee z)$ and $(\mathrm{nor}_{3c})(\langle x,y,z\rangle) = \neg(\neg x \vee \neg y \vee \neg z)$.

(19)[2] For all elements x, y, z of *Boolean* holds $\mathrm{and}_3(\langle x,y,z\rangle) = (\mathrm{nor}_{3c})(\langle x,y,z\rangle)$ and $(\mathrm{and}_{3a})(\langle x,y,z\rangle) = (\mathrm{nor}_{3b})(\langle z,y,x\rangle)$ and $(\mathrm{and}_{3b})(\langle x,y,z\rangle) = (\mathrm{nor}_{3a})(\langle z,y,x\rangle)$ and $(\mathrm{and}_{3c})(\langle x,y,z\rangle) = \mathrm{nor}_3(\langle x,y,z\rangle)$.

(20) For all elements x, y, z of *Boolean* holds $\mathrm{or}_3(\langle x,y,z\rangle) = (\mathrm{nand}_{3c})(\langle x,y,z\rangle)$ and $(\mathrm{or}_{3a})(\langle x,y,z\rangle) = (\mathrm{nand}_{3b})(\langle z,y,x\rangle)$ and $(\mathrm{or}_{3b})(\langle x,y,z\rangle) = (\mathrm{nand}_{3a})(\langle z,y,x\rangle)$ and $(\mathrm{or}_{3c})(\langle x,y,z\rangle) = \mathrm{nand}_3(\langle x,y,z\rangle)$.

(21) $\mathrm{and}_3(\langle 0,0,0\rangle) = 0$ and $\mathrm{and}_3(\langle 0,0,1\rangle) = 0$ and $\mathrm{and}_3(\langle 0,1,0\rangle) = 0$ and $\mathrm{and}_3(\langle 0,1,1\rangle) = 0$ and $\mathrm{and}_3(\langle 1,0,0\rangle) = 0$ and $\mathrm{and}_3(\langle 1,0,1\rangle) = 0$ and $\mathrm{and}_3(\langle 1,1,0\rangle) = 0$ and $\mathrm{and}_3(\langle 1,1,1\rangle) = 1$.

(22) $(\mathrm{and}_{3a})(\langle 0,0,0\rangle) = 0$ and $(\mathrm{and}_{3a})(\langle 0,0,1\rangle) = 0$ and $(\mathrm{and}_{3a})(\langle 0,1,0\rangle) = 0$ and $(\mathrm{and}_{3a})(\langle 0,1,1\rangle) = 1$ and $(\mathrm{and}_{3a})(\langle 1,0,0\rangle) = 0$ and $(\mathrm{and}_{3a})(\langle 1,0,1\rangle) = 0$ and $(\mathrm{and}_{3a})(\langle 1,1,0\rangle) = 0$ and $(\mathrm{and}_{3a})(\langle 1,1,1\rangle) = 0$.

(23) $(\mathrm{and}_{3b})(\langle 0,0,0\rangle) = 0$ and $(\mathrm{and}_{3b})(\langle 0,0,1\rangle) = 1$ and $(\mathrm{and}_{3b})(\langle 0,1,0\rangle) = 0$ and $(\mathrm{and}_{3b})(\langle 0,1,1\rangle) = 0$ and $(\mathrm{and}_{3b})(\langle 1,0,0\rangle) = 0$ and $(\mathrm{and}_{3b})(\langle 1,0,1\rangle) = 0$ and $(\mathrm{and}_{3b})(\langle 1,1,0\rangle) = 0$ and $(\mathrm{and}_{3b})(\langle 1,1,1\rangle) = 0$.

(24) $(\mathrm{and}_{3c})(\langle 0,0,0\rangle) = 1$ and $(\mathrm{and}_{3c})(\langle 0,0,1\rangle) = 0$ and $(\mathrm{and}_{3c})(\langle 0,1,0\rangle) = 0$ and $(\mathrm{and}_{3c})(\langle 0,1,1\rangle) = 0$ and $(\mathrm{and}_{3c})(\langle 1,0,0\rangle) = 0$ and $(\mathrm{and}_{3c})(\langle 1,0,1\rangle) = 0$ and $(\mathrm{and}_{3c})(\langle 1,1,0\rangle) = 0$ and $(\mathrm{and}_{3c})(\langle 1,1,1\rangle) = 0$.

(25) $\mathrm{or}_3(\langle 0,0,0\rangle) = 0$ and $\mathrm{or}_3(\langle 0,0,1\rangle) = 1$ and $\mathrm{or}_3(\langle 0,1,0\rangle) = 1$ and $\mathrm{or}_3(\langle 0,1,1\rangle) = 1$ and $\mathrm{or}_3(\langle 1,0,0\rangle) = 1$ and $\mathrm{or}_3(\langle 1,0,1\rangle) = 1$ and $\mathrm{or}_3(\langle 1,1,0\rangle) = 1$ and $\mathrm{or}_3(\langle 1,1,1\rangle) = 1$.

(26) $(\mathrm{or}_{3a})(\langle 0,0,0\rangle) = 1$ and $(\mathrm{or}_{3a})(\langle 0,0,1\rangle) = 1$ and $(\mathrm{or}_{3a})(\langle 0,1,0\rangle) = 1$ and $(\mathrm{or}_{3a})(\langle 0,1,1\rangle) = 1$ and $(\mathrm{or}_{3a})(\langle 1,0,0\rangle) = 0$ and $(\mathrm{or}_{3a})(\langle 1,0,1\rangle) = 1$ and $(\mathrm{or}_{3a})(\langle 1,1,0\rangle) = 1$ and $(\mathrm{or}_{3a})(\langle 1,1,1\rangle) = 1$.

(27) $(\mathrm{or}_{3b})(\langle 0,0,0\rangle) = 1$ and $(\mathrm{or}_{3b})(\langle 0,0,1\rangle) = 1$ and $(\mathrm{or}_{3b})(\langle 0,1,0\rangle) = 1$ and $(\mathrm{or}_{3b})(\langle 0,1,1\rangle) = 1$ and $(\mathrm{or}_{3b})(\langle 1,0,0\rangle) = 1$ and $(\mathrm{or}_{3b})(\langle 1,0,1\rangle) = 1$ and $(\mathrm{or}_{3b})(\langle 1,1,0\rangle) = 0$ and $(\mathrm{or}_{3b})(\langle 1,1,1\rangle) = 1$.

(28) $(\mathrm{or}_{3c})(\langle 0,0,0\rangle) = 1$ and $(\mathrm{or}_{3c})(\langle 0,0,1\rangle) = 1$ and $(\mathrm{or}_{3c})(\langle 0,1,0\rangle) = 1$ and $(\mathrm{or}_{3c})(\langle 0,1,1\rangle) = 1$ and $(\mathrm{or}_{3c})(\langle 1,0,0\rangle) = 1$ and $(\mathrm{or}_{3c})(\langle 1,0,1\rangle) = 1$ and $(\mathrm{or}_{3c})(\langle 1,1,0\rangle) = 1$ and $(\mathrm{or}_{3c})(\langle 1,1,1\rangle) = 0$.

(29) $\mathrm{xor}_3(\langle 0,0,0\rangle) = 0$ and $\mathrm{xor}_3(\langle 0,0,1\rangle) = 1$ and $\mathrm{xor}_3(\langle 0,1,0\rangle) = 1$ and $\mathrm{xor}_3(\langle 0,1,1\rangle) = 0$ and $\mathrm{xor}_3(\langle 1,0,0\rangle) = 1$ and $\mathrm{xor}_3(\langle 1,0,1\rangle) = 0$ and $\mathrm{xor}_3(\langle 1,1,0\rangle) = 0$ and $\mathrm{xor}_3(\langle 1,1,1\rangle) = 1$.

[2] The proposition (18) has been removed.

2. 2'S COMPLEMENT CIRCUIT PROPERTIES

Let x, b be sets. The functor $\mathrm{CompStr}(x,b)$ yielding an unsplit non void strict non empty many sorted signature with arity held in gates and Boolean denotation held in gates is defined as follows:

(Def. 33) $\mathrm{CompStr}(x,b) = \mathrm{1GateCircStr}(\langle x,b\rangle, \mathrm{xor}_{2a})$.

Let x, b be sets. The functor $\mathrm{CompCirc}(x,b)$ yielding a strict Boolean circuit of $\mathrm{CompStr}(x,b)$ with denotation held in gates is defined by:

(Def. 34) $\mathrm{CompCirc}(x,b) = \mathrm{1GateCircuit}(x,b,\mathrm{xor}_{2a})$.

Let x, b be sets. The functor $\mathrm{CompOutput}(x,b)$ yielding an element of $\mathrm{InnerVertices}(\mathrm{CompStr}(x,b))$ is defined by:

(Def. 35) $\mathrm{CompOutput}(x,b) = \langle\langle x,b\rangle, \mathrm{xor}_{2a}\rangle$.

Let x, b be sets. The functor $\mathrm{IncrementStr}(x,b)$ yielding an unsplit non void strict non empty many sorted signature with arity held in gates and Boolean denotation held in gates is defined as follows:

(Def. 36) $\mathrm{IncrementStr}(x,b) = \mathrm{1GateCircStr}(\langle x,b\rangle, \mathrm{and}_{2a})$.

Let x, b be sets. The functor $\mathrm{IncrementCirc}(x,b)$ yields a strict Boolean circuit of $\mathrm{IncrementStr}(x,b)$ with denotation held in gates and is defined as follows:

(Def. 37) $\mathrm{IncrementCirc}(x,b) = \mathrm{1GateCircuit}(x,b,\mathrm{and}_{2a})$.

Let x, b be sets. The functor $\mathrm{IncrementOutput}(x,b)$ yielding an element of $\mathrm{InnerVertices}(\mathrm{IncrementStr}(x,b))$ is defined by:

(Def. 38) $\mathrm{IncrementOutput}(x,b) = \langle\langle x,b\rangle, \mathrm{and}_{2a}\rangle$.

Let x, b be sets. The functor $\mathrm{BitCompStr}(x,b)$ yields an unsplit non void strict non empty many sorted signature with arity held in gates and Boolean denotation held in gates and is defined by:

(Def. 39) $\mathrm{BitCompStr}(x,b) = \mathrm{CompStr}(x,b) +\cdot \mathrm{IncrementStr}(x,b)$.

Let x, b be sets. The functor $\mathrm{BitCompCirc}(x,b)$ yields a strict Boolean circuit of $\mathrm{BitCompStr}(x,b)$ with denotation held in gates and is defined as follows:

(Def. 40) $\mathrm{BitCompCirc}(x,b) = \mathrm{CompCirc}(x,b) +\cdot \mathrm{IncrementCirc}(x,b)$.

One can prove the following propositions:

(30) For all non pair sets x, b holds $\mathrm{InnerVertices}(\mathrm{CompStr}(x,b))$ is a binary relation.

(31) For all non pair sets x, b holds $x \in$ the carrier of $\mathrm{CompStr}(x,b)$ and $b \in$ the carrier of $\mathrm{CompStr}(x,b)$ and $\langle\langle x,b\rangle, \mathrm{xor}_{2a}\rangle \in$ the carrier of $\mathrm{CompStr}(x,b)$.

(32) For all non pair sets x, b holds the carrier of $\mathrm{CompStr}(x,b) = \{x,b\} \cup \{\langle\langle x,b\rangle, \mathrm{xor}_{2a}\rangle\}$.

(33) For all non pair sets x, b holds $\mathrm{InnerVertices}(\mathrm{CompStr}(x,b)) = \{\langle\langle x,b\rangle, \mathrm{xor}_{2a}\rangle\}$.

(34) For all non pair sets x, b holds $\langle\langle x,b\rangle, \mathrm{xor}_{2a}\rangle \in \mathrm{InnerVertices}(\mathrm{CompStr}(x,b))$.

(35) For all non pair sets x, b holds $\mathrm{InputVertices}(\mathrm{CompStr}(x,b)) = \{x,b\}$.

(36) For all non pair sets x, b holds $x \in \mathrm{InputVertices}(\mathrm{CompStr}(x,b))$ and $b \in \mathrm{InputVertices}(\mathrm{CompStr}(x,b))$.

(37) For all non pair sets x, b holds $\mathrm{InputVertices}(\mathrm{CompStr}(x,b))$ has no pairs.

(38) For all non pair sets x, b holds $\mathrm{InnerVertices}(\mathrm{IncrementStr}(x,b))$ is a binary relation.

(39) For all non pair sets x, b holds $x \in$ the carrier of IncrementStr(x,b) and $b \in$ the carrier of IncrementStr(x,b) and $\langle\langle x,b\rangle, \text{and}_{2a}\rangle \in$ the carrier of IncrementStr(x,b).

(40) For all non pair sets x, b holds the carrier of IncrementStr$(x,b) = \{x,b\} \cup \{\langle\langle x,b\rangle, \text{and}_{2a}\rangle\}$.

(41) For all non pair sets x, b holds InnerVertices(IncrementStr$(x,b)) = \{\langle\langle x,b\rangle, \text{and}_{2a}\rangle\}$.

(42) For all non pair sets x, b holds $\langle\langle x,b\rangle, \text{and}_{2a}\rangle \in$ InnerVertices(IncrementStr$(x,b))$.

(43) For all non pair sets x, b holds InputVertices(IncrementStr$(x,b)) = \{x,b\}$.

(44) For all non pair sets x, b holds $x \in$ InputVertices(IncrementStr$(x,b))$ and $b \in$ InputVertices(IncrementStr$(x,b))$.

(45) For all non pair sets x, b holds InputVertices(IncrementStr$(x,b))$ has no pairs.

(46) For all non pair sets x, b holds InnerVertices(BitCompStr$(x,b))$ is a binary relation.

(47) Let x, b be non pair sets. Then

(i) $x \in$ the carrier of BitCompStr(x,b),

(ii) $b \in$ the carrier of BitCompStr(x,b),

(iii) $\langle\langle x,b\rangle, \text{xor}_{2a}\rangle \in$ the carrier of BitCompStr(x,b), and

(iv) $\langle\langle x,b\rangle, \text{and}_{2a}\rangle \in$ the carrier of BitCompStr(x,b).

(48) For all non pair sets x, b holds the carrier of BitCompStr$(x,b) = \{x,b\} \cup \{\langle\langle x,b\rangle, \text{xor}_{2a}\rangle, \langle\langle x,b\rangle, \text{and}_{2a}\rangle\}$.

(49) For all non pair sets x, b holds InnerVertices(BitCompStr$(x,b)) = \{\langle\langle x,b\rangle, \text{xor}_{2a}\rangle, \langle\langle x,b\rangle, \text{and}_{2a}\rangle\}$.

(50) For all non pair sets x, b holds $\langle\langle x,b\rangle, \text{xor}_{2a}\rangle \in$ InnerVertices(BitCompStr$(x,b))$ and $\langle\langle x, b\rangle, \text{and}_{2a}\rangle \in$ InnerVertices(BitCompStr$(x,b))$.

(51) For all non pair sets x, b holds InputVertices(BitCompStr$(x,b)) = \{x,b\}$.

(52) For all non pair sets x, b holds $x \in$ InputVertices(BitCompStr$(x,b))$ and $b \in$ InputVertices(BitCompStr$(x,b))$.

(53) For all non pair sets x, b holds InputVertices(BitCompStr$(x,b))$ has no pairs.

(54) For all non pair sets x, b and for every state s of CompCirc(x,b) holds (Following(s))(CompOutput$(x,b)) = (\text{xor}_{2a})(\langle s(x), s(b)\rangle)$ and (Following(s))$(x) = s(x)$ and (Following(s))$(b) = s(b)$.

(55) Let x, b be non pair sets, s be a state of CompCirc(x,b), and a_1, a_2 be elements of *Boolean*. If $a_1 = s(x)$ and $a_2 = s(b)$, then (Following(s))(CompOutput$(x,b)) = \neg a_1 \oplus a_2$ and (Following(s))$(x) = a_1$ and (Following(s))$(b) = a_2$.

(56) For all non pair sets x, b and for every state s of BitCompCirc(x,b) holds (Following(s))(CompOutput$(x,b)) = (\text{xor}_{2a})(\langle s(x), s(b)\rangle)$ and (Following(s))$(x) = s(x)$ and (Following(s))$(b) = s(b)$.

(57) Let x, b be non pair sets, s be a state of BitCompCirc(x,b), and a_1, a_2 be elements of *Boolean*. If $a_1 = s(x)$ and $a_2 = s(b)$, then (Following(s))(CompOutput$(x,b)) = \neg a_1 \oplus a_2$ and (Following(s))$(x) = a_1$ and (Following(s))$(b) = a_2$.

(58) For all non pair sets x, b and for every state s of IncrementCirc(x,b) holds (Following(s))(IncrementOutput$(x,b)) = (\text{and}_{2a})(\langle s(x), s(b)\rangle)$ and (Following(s))$(x) = s(x)$ and (Following(s))$(b) = s(b)$.

(59) Let x, b be non pair sets, s be a state of IncrementCirc(x,b), and a_1, a_2 be elements of *Boolean*. If $a_1 = s(x)$ and $a_2 = s(b)$, then $(\mathrm{Following}(s))(\mathrm{IncrementOutput}(x,b)) = \neg a_1 \wedge a_2$ and $(\mathrm{Following}(s))(x) = a_1$ and $(\mathrm{Following}(s))(b) = a_2$.

(60) For all non pair sets x, b and for every state s of BitCompCirc(x,b) holds $(\mathrm{Following}(s))(\mathrm{IncrementOutput}(x,b)) = (\mathrm{and}_{2a})(\langle s(x), s(b)\rangle)$ and $(\mathrm{Following}(s))(x) = s(x)$ and $(\mathrm{Following}(s))(b) = s(b)$.

(61) Let x, b be non pair sets, s be a state of BitCompCirc(x,b), and a_1, a_2 be elements of *Boolean*. If $a_1 = s(x)$ and $a_2 = s(b)$, then $(\mathrm{Following}(s))(\mathrm{IncrementOutput}(x,b)) = \neg a_1 \wedge a_2$ and $(\mathrm{Following}(s))(x) = a_1$ and $(\mathrm{Following}(s))(b) = a_2$.

(62) Let x, b be non pair sets and s be a state of BitCompCirc(x,b). Then $(\mathrm{Following}(s))(\mathrm{CompOutput}(x,b)) = (\mathrm{xor}_{2a})(\langle s(x), s(b)\rangle)$ and $(\mathrm{Following}(s))(\mathrm{IncrementOutput}(x,b)) = (\mathrm{and}_{2a})(\langle s(x), s(b)\rangle)$ and $(\mathrm{Following}(s))(x) = s(x)$ and $(\mathrm{Following}(s))(b) = s(b)$.

(63) Let x, b be non pair sets, s be a state of BitCompCirc(x,b), and a_1, a_2 be elements of *Boolean*. Suppose $a_1 = s(x)$ and $a_2 = s(b)$. Then $(\mathrm{Following}(s))(\mathrm{CompOutput}(x,b)) = \neg a_1 \oplus a_2$ and $(\mathrm{Following}(s))(\mathrm{IncrementOutput}(x,b)) = \neg a_1 \wedge a_2$ and $(\mathrm{Following}(s))(x) = a_1$ and $(\mathrm{Following}(s))(b) = a_2$.

(64) For all non pair sets x, b and for every state s of BitCompCirc(x,b) holds Following(s) is stable.

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